Towards an efficient implementation of CADNA in the BLAS: Example of DgemmCADNA routine

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Summary

1. Motivations

2. The CADNA Library

3. CADNA Implementation in scientific libraries

4. Conclusion
1. Motivations

2. The CADNA Library

3. CADNA Implementation in scientific libraries

4. Conclusion
Numerical simulation

Several approximations!

Computed results can be wrong!

- Round-off errors at each elementary arithmetic operation
- Detect and control these errors
  - Numerical validation
Numerical validation : Tools/Methods (1)

Methods for accurate computations

- Multiple precision arithmetic :
  - ex : MPFR, Gnu MP, MPFI.

- Compensated methods :
  - compensated summation algorithms, compensated dot product algorithms...
Numerical validation : Tools/Methods (2)

**Methods for rounding error analysis**

- Inverse analysis :
  - provides error bounds for the computed results.

- Interval arithmetic : the result of an arithmetic operation between two intervals contains all values that can be obtained by performing this operation on elements from each interval.

- Probabilistic approach :
  - uses a random rounding mode (CESTAC Method);
  - estimates the number of exact significant digits of any computed result.
Find an adapted tool for the industrial context (EDF)

The CADNA Library [SJDC07].
2. The CADNA Library

1. Motivations

2. The CADNA Library

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4. Conclusion
The CESTAC method:

The CESTAC method (Contrôle et Estimation Stochastique des Arrondis de Calculs) was proposed by M. La Porte and J. Vignes in 1974 [VLP74].

It consists in running the same code several times with different round-off error propagations. Then, different results are obtained.

- the part that is common to all the different results is assumed to be also in common with the mathematical result;
- the part that is different in the results is affected by the round-off errors.
Discrete Stochastic Arithmetic

- $N$ different runs with random rounding mode ($+\infty$ ou $-\infty$ with the probability 0.5);
- $N$ different results $R_i$:
  - choosing as the computed result the mean value $\bar{R}$ of $R_i$;
  - estimating $C_R$ the number of exact significant decimal digits of $\bar{R}$.
- $N = 3$
  - $X = (X_1, X_2, X_3)$
  - $\forall \Omega \in (+, -, \times, /), X\Omega Y = (X_1\omega Y_1, X_2\omega Y_2, X_3\omega Y_3)$

- If $C_R \leq 0$ or $\forall i, R_i = 0$, a result $R$ is a computed zero (@.0).
- New order relationships.

Discrete Stochastic Arithmetic (DSA).
The CADNA library

The CADNA library implements Discrete Stochastic Arithmetic. It allows the estimation of round-off error propagation in any scientific program [JCL10].

More precisely, CADNA enables one to:
- estimate the numerical quality of any result
- control branching statements
- perform a dynamic numerical debugging
- take into account uncertainty on data.

CADNA is a library which can be used with Fortran or C++ programs and also with MPI parallel programs. CADNA can be downloaded from http://www.lip6.fr/cadna
How to use the Cadna Library

- CADNA provides two new numerical types, the stochastic types (3 floating point variables x, y, z and a hidden variable acc):
  - type *(single_st)* in single precision
  - type *(double_st)* in double precision.

- All the operators and mathematical functions are overloaded for these types.

- To use the library:
  1. declaration of the CADNA library
  2. initialization of the CADNA library
  3. substitution of the floating point type by stochastic types
  4. change of output statements to print stochastic results with their accuracy
  5. termination of the CADNA library
### High Performance Computing at EDF R&D

<table>
<thead>
<tr>
<th>Codes</th>
<th>Tools</th>
<th>Hardware</th>
</tr>
</thead>
<tbody>
<tr>
<td>Code_Aster</td>
<td>MPI/OpenMP</td>
<td>Ivanoe</td>
</tr>
<tr>
<td>Code_Saturne</td>
<td>BLAS/LAPACK</td>
<td>Blue Gene</td>
</tr>
<tr>
<td>TELEMAC</td>
<td>MUMPS/PASTIX</td>
<td>Clamart2</td>
</tr>
<tr>
<td>...</td>
<td>...</td>
<td>Z600</td>
</tr>
</tbody>
</table>

Is it possible to study the numerical quality of every industrial code with CADNA?
3. CADNA Implementation in scientific libraries

1. Motivations

2. The CADNA Library

3. CADNA Implementation in scientific libraries
   - The communication standards MPI and BLACS
   - CADNA implementation in BLAS routines

4. Conclusion
Different extensions for CADNA

1. MPI extension for CADNA : CADNA MPI

2. BLACS extension for CADNA : CADNA BLACS

3. Efficient implementation of CADNA in BLAS
**MPI/BLACS extensions for CADNA**

- Definition of stochastic types to exchange data
- Definition of reduction operators
- C/C++ (MPI2), Fortran 90 (MPI1)
  - (--) The sending time of a stochastic float is 4 times more long than a normal float one.
    - Size of stochastic type = 4 times size of normal float
  - (++) It is possible to use CADNA with any code using MPI and BLACS.
BLAS : Basic Linear Algebra Subprograms

Functionality
- Level 1 : vectors operations (ex xAXPY);
- Level 2 : matrix-vectors operations (ex xGEMV);
- Level 3 : matrix-matrix operations (ex xGEMM).

Implementations

<table>
<thead>
<tr>
<th>versions</th>
<th>daxpy</th>
<th>dgemv</th>
<th>dgemm</th>
</tr>
</thead>
<tbody>
<tr>
<td>Netlib</td>
<td>1.18482</td>
<td>1.15347</td>
<td>1.35378</td>
</tr>
<tr>
<td>Mkl 1 threads</td>
<td>2.02116</td>
<td>2.11232</td>
<td>7.53686</td>
</tr>
<tr>
<td>Goto 1 threads</td>
<td>2.86331</td>
<td>2.12331</td>
<td>7.52166</td>
</tr>
<tr>
<td>Mkl 8 threads</td>
<td>1.63618</td>
<td>2.79974</td>
<td>58.0523</td>
</tr>
<tr>
<td>Goto 8 threads</td>
<td>1.63618</td>
<td>4.60287</td>
<td>56.3343</td>
</tr>
</tbody>
</table>

Table: GFLOPS for daxpy, dgemv et dgemm : 4096*4096 matrix (4096 vector).
How to use CADNA with Blas routines?

The easiest solution (V1) : Remplaced \textit{float} by \textit{float\_st} et \textit{double} by \textit{double\_st} :

\begin{verbatim}
void cblas_dgemm(const enum CBLAS_ORDER Order, const enum CBLAS_TRANSPOSE TransA, const enum CBLAS_TRANSPOSE TransB, const int M, const int N, const int K, const double\_st alpha, const double\_st *A, const int lda, const double\_st *B, const int ldb, const double\_st beta, double\_st *C, const int ldc);
\end{verbatim}

Linalg : A template version of BLAS ; it can be used with stochastic types
Direct Implementations of DGEMM with CADNA

**Figure**: Versions with and without CADNA
Direct Implementations of DGEMM with CADNA(2)

Figure: Overhead due to the CESTAC Method
Why these overheads?

- An overhead greater than 1000 for a 1024*1024 matrix
  - DGEMM with 3 inner loops => cache misses
  - Use of stochastic types and the discrete stochastic arithmetic
  - Random rounding mode \( V1 \ > \ 7xV2 \)
Efficient implementation of DgemmCADNA

« Implementing matrix multiplication so that near-optimal performance is attained requires a thorough understanding of how the operation must be layered at the macro level in combination with careful engineering of high-performance kernels at the micro level. »

K. Goto, 2008 [GVDG08].

Solutions to reduce the overhead:

1. Efficient use of data (memory access)
2. Minimize the CESTAC Method impact
3. Optimize the inner loop
Efficient use of data or data reuse

- Use tiled algorithms
- Optimize cache locality
- Exploit temporal and spacial locality
- Reduce cache misses
- Reduce TLB misses
An Iterative tiled algorithm

\[
\begin{bmatrix}
C_{11} & C_{12} & \ldots & C_{1N} \\
C_{21} & C_{22} & \ldots & C_{2N} \\
\vdots & \vdots & \ddots & \vdots \\
C_{N1} & C_{N2} & \ldots & C_{NN}
\end{bmatrix} =
\begin{bmatrix}
A_{11} & A_{12} & \ldots & A_{1N} \\
A_{21} & A_{22} & \ldots & A_{2N} \\
\vdots & \vdots & \ddots & \vdots \\
A_{N1} & A_{N2} & \ldots & A_{NN}
\end{bmatrix}
\times
\begin{bmatrix}
B_{11} & B_{12} & \ldots & B_{1N} \\
B_{21} & B_{22} & \ldots & B_{2N} \\
\vdots & \vdots & \ddots & \vdots \\
B_{N1} & B_{N2} & \ldots & B_{NN}
\end{bmatrix}
\]

every \( C_{ij} \) is computed by:

\[
C_{ij} = \sum_{k=1}^{N} A_{ik} B_{kj}
\]
A recursive tiled algorithm DGBRn

\[
\begin{bmatrix}
C_{00} & C_{01} \\
C_{10} & C_{11}
\end{bmatrix}
= 
\begin{bmatrix}
A_{00} & A_{01} \\
A_{10} & A_{11}
\end{bmatrix}
\times 
\begin{bmatrix}
B_{00} & B_{01} \\
B_{10} & B_{11}
\end{bmatrix}
\] (1)
An iterative tiled algorithm based on the hardware (hierarchical memory) DGBIn

3 levels of partitioning. One level for every cache level. The matrix (submatrices) is partitioned in submatrices (blocks). At each step, 3 blocks must fit in this level of cache memory.

1. First level for Cache L3
2. Second level for Cache L2
3. Third level for Cache L1

\[ A(n \times n) = \begin{bmatrix} A_{11} & A_{12} & \ldots & A_{1N} \\ A_{21} & A_{22} & \ldots & A_{2N} \\ \vdots & \vdots & \ddots & \vdots \\ A_{N1} & A_{N2} & \ldots & A_{NN} \end{bmatrix} \]
**BDL : Block Data Layout**

- **Column Major order**
  \[
  M = \begin{bmatrix}
  1 & 2 & 3 \\
  4 & 5 & 6 \\
  7 & 8 & 9 \\
  \end{bmatrix}
  \Rightarrow \begin{bmatrix}
  1 & 4 & 7 \\
  2 & 5 & 8 \\
  3 & 6 & 9 \\
  \end{bmatrix}
  \]

- **Row Major order**
  \[
  M = \begin{bmatrix}
  1 & 2 & 3 \\
  4 & 5 & 6 \\
  7 & 8 & 9 \\
  \end{bmatrix}
  \Rightarrow \begin{bmatrix}
  1 & 2 & 3 \\
  4 & 5 & 6 \\
  7 & 8 & 9 \\
  \end{bmatrix}
  \]
BDL : *Block Data Layout* (2)

Consider matrix $A(n \times n)$ partitioned in $N \times N$ submatrices $A_{ij}$:

$$
A(n \times n) = \begin{bmatrix}
A_{11} & A_{12} & \cdots & A_{1N} \\
A_{21} & A_{22} & \cdots & A_{2N} \\
\vdots & \vdots & \ddots & \vdots \\
A_{N1} & A_{N2} & \cdots & A_{NN}
\end{bmatrix}
$$

$$
A_{ij}(p \times p) = \begin{bmatrix}
a_{11} & a_{12} & \cdots & a_{1p} \\
a_{21} & a_{22} & \cdots & a_{2p} \\
\vdots & \vdots & \ddots & \vdots \\
a_{p1} & a_{p2} & \cdots & a_{pp}
\end{bmatrix}
$$

Data within one such block $A_{ij}$ are mapped onto contiguous memory:

$$
\begin{bmatrix}
a_{11} & a_{12} & \cdots & a_{1p} & a_{21} & a_{22} & \cdots & a_{2p} & \cdots & a_{p1} & a_{p2} & \cdots & a_{pp}
\end{bmatrix}
$$

Theses blocks are arranged in *row-major order*:

$$
\begin{bmatrix}
A_{11} & A_{12} & \cdots & A_{1N} & A_{21} & A_{22} & \cdots & A_{2N} & \cdots & A_{N1} & A_{N2} & \cdots & A_{NN}
\end{bmatrix}
$$
Reduce the impact of DSA

Unroll every Cadna arithmetic operation: NO MORE operator overloading

\[ C[i] = A[i] + B[i] ; \]

\[ C[i].x = A[i].x + B[i].x ; \]
\[ \text{if (random) } \text{rnd_switch}() ; \]
\[ C[i].y = A[i].y + B[i].y ; \]
\[ \text{if (random) } \text{rnd_switch}() ; \]
\[ C[i].z = A[i].z + B[i].z ; \]
\[ \text{rnd_switch}() ; \]
Reduce the impact of DSA (2)

less calls to \texttt{rnd\_switch()}

\begin{verbatim}
\textbf{if (random) \texttt{rnd\_switch()}}
\end{verbatim}
\begin{verbatim}
C[i].x = A[i].x + B[i].x ;
C[i].z = A[i].z + B[i].z ;
C[i+1].z = A[i+1].z + B[i+1].z ;
C[i+2].x = A[i+2].x + B[i+2].x ;
C[i+2].y = A[i+2].y + B[i+2].y ;
C[i+3].x = A[i+3].x + B[i+3].x ;
C[i+3].y = A[i+3].y + B[i+3].y ;
\textbf{rnd\_switch();}
\end{verbatim}
\begin{verbatim}
C[i].y = A[i].y + B[i].y ;
C[i+1].x = A[i+1].x + B[i+1].x ;
C[i+1].y = A[i+1].y + B[i+1].y ;
C[i+2].z = A[i+2].z + B[i+2].z ;
C[i+3].y = A[i+3].y + B[i+3].y ;
C[i+3].z = A[i+3].z + B[i+3].z ;
\end{verbatim}
Optimize the kernel

Listing 1 – Inner loops

```c
for(int i = 0; i < nb_block; i++){
    for(int k = 0; k < nb_block; k++){
        for(int j = 0; j < nb_block; j++){
            Cij = Aik * Bkj /*kernel*/
        }
    }
}
```

\[
\begin{align*}
C_{00} &= A_{00} \times B_{00} \\
C_{01} &= A_{00} \times B_{01} \\
C_{00} &= A_{01} \times B_{10} \\
C_{01} &= A_{01} \times B_{11} \\
C_{10} &= A_{10} \times B_{00} \\
C_{11} &= A_{10} \times B_{01} \\
C_{10} &= A_{11} \times B_{10} \\
C_{11} &= A_{11} \times B_{11}
\end{align*}
\]
Results

**Figure:** Different versions of DgemmCADNA
Results (2)

**Figure**: Comparison to GotoBLAS
4. Conclusion

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Conclusion et Future work

**Conclusion**

- CADNA extensions for MPI and BLACS;
- DgemmCADNA subroutine:
  - 45x faster than the first version
  - Gotoblas 25x faster than DgemmCADNA

**Future work**

- Include the CADNA autovalidation;
- Theorical proof of the CESTAC Method modification
- Work on the other blas routines (level 1 and level 2)
- Experimental test phase for the implemented routines in a industrial codes (TELEMAC)
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Thanks!

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